

Packing Chromatic Number of Cycle Related Graphs

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Abstract

The packing chromatic number $\chi_\rho(G)$ of a graph G is the smallest integer k for which there exists a mapping $\pi : V(G) \rightarrow \{1, 2, \dots, k\}$ such that any two vertices of color i are at distance at least $i + 1$. It is a frequency assignment problem used in wireless networks, which is also called Broadcast coloring. It is proved that packing coloring is NP-complete for general graphs and even for trees. In this paper, we give the packing chromatic number for some classes of cycles.

Keywords: Packing chromatic number, kC_n - linear, barbell graph, split graph, total graph, crown graph, lollipop, tadpole.

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1 Introduction

Let G be a connected graph and k be an integer, $k \geq 1$. A packing k -coloring of a graph G is a mapping $\pi : V(G) \rightarrow \{1, 2, \dots, k\}$ such that any two vertices of color i are at distance at least $i + 1$. Thus, the vertices of G are partitioned into different color classes X_1, X_2, \dots, X_k , where every X_i is an i -packing of G . The i -packing number of G , denoted by $\rho_i(G)$, is the maximum cardinality of an i -packing that occurs in G . The packing chromatic number $\chi_\rho(G)$ of G is the smallest integer k for which G has packing k -coloring. The concept of packing coloring emerges from the area of frequency assignment in wireless networks and was introduced by Goddard et al. [8] by the name Broadcast coloring. It has several applications such as resource replacement and biological diversity. The term packing chromatic number was introduced by Bresar [3].

Goddard et al. [8] proved that the packing coloring problem is NP-complete for general graphs and Fiala and Golovach [6] proved that it is NP-complete even for trees. It is proved that packing coloring problem is solvable in polynomial time for graphs whose treewidth and diameter are both bounded [6] and for cographs and split graphs [8]. Sloper [14] studied a special type of packing coloring, called eccentric chromatic coloring and proved that the infinite 3-regular tree has packing chromatic number 7. For the infinite planar square lattice \mathbb{Z}^2 , $10 \leq \chi_\rho(\mathbb{Z}^2) \leq 17$ [5, 9]. The packing coloring of distance graphs were studied in [4, 15]. For the infinite hexagonal lattice \mathbb{H} , $\chi_\rho(\mathbb{H}) = 7$ [3].

Argiroffo et al. [1, 2] proved that the packing coloring is solvable in polynomial time for the class of $(q, q - 4)$ graphs, partner limited graphs and for an infinite subclass of lobsters,

including caterpillars. In [7, 13] it is proved that the infinite, planar triangular lattice and the three dimensional square lattice have unbounded packing chromatic number. In this paper, we study the packing chromatic number of kC_n - linear, barbell graph, split graph, total graph, crown graph, lollipop and tadpole.

2 Main Results

Lemma 2.1. [8] Let H be a subgraph of G . Then $\chi_\rho(H) \leq \chi_\rho(G)$.

Lemma 2.2. [8] $\chi_\rho(C_n) = \begin{cases} 3 & \text{when } n \text{ is a multiple of } 4 \\ 4 & \text{when } n \text{ is not a multiple of } 4. \end{cases}$

Lemma 2.3. [8] Let G be a complete graph on n vertices. Then $\chi_\rho(G) = n$.

Definition 2.4. A graph is called a kC_{2n} - linear if it is a connected graph with k blocks, each of which is isomorphic to C_{2n} . Moreover, if u and v are cut-vertices in the same block B_i , then $d(u, v) = n$.

Theorem 2.5. For $kC_{2n}, n \geq 4, \chi_\rho(kC_{2n}) = 3$ when $n \equiv 0 \pmod 4$.

Proof: Case 1: $n = 8r$.

We begin with the cut-vertex common to B_i, B_{i+1} and label it as 2. Label every 4-sequence of consecutive vertices on B_i in the clockwise sense as 2,1,3,1 beginning from the cut-vertex. See Figure 1(a).

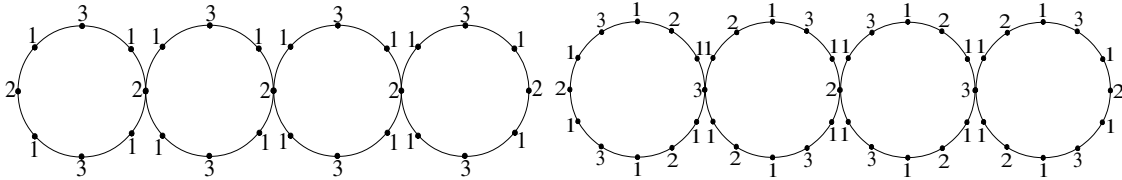


Figure 1: (a) $\chi_\rho(4C_8) = 3$ (b) $\chi_\rho(4C_{12}) = 3$.

Case 2: $n = 8r + 4$.

Label the cut-vertex common to B_i, B_{i+1} as 3 or 2 according as i is odd or even respectively. Label every 4-sequence of consecutive vertices on B_i in the clockwise sense as 3,1,2,1 or 2,1,3,1 beginning from the cut-vertex with label 3 or 2. See Figure 1(b). By Lemma 2.1 and 2.2, $\chi_\rho(kC_{2n}) = 3$. ■

Definition 2.6. [11] A gear graph is obtained from the wheel W_n by adding a vertex between every pair of adjacent vertices of the n -cycle and it is denoted by G_n , where n is the number of vertices of wheel graph. G_n has $2n + 1$ vertices and $3n$ edges.

Theorem 2.7. For the gear graph $G_n, \chi_\rho(G_n) = 4$, when n is even.

Proof: Let u and v be the vertices of G_n such that $deg(u) = n$ and $deg(v) = 3$. Color the vertices in the outercycle with the sequence 1213, \dots starting from the vertex v in any direction and fix color 4 to the vertex u . Thus, $\chi_\rho(G_n) \leq 4$.

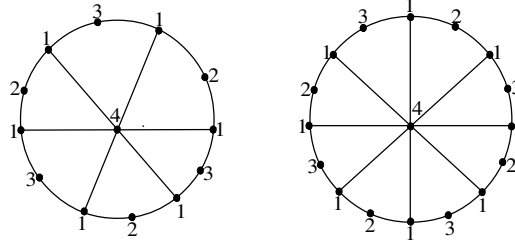


Figure 2 : $\chi_\rho(G_6) = 4$ and $\chi_\rho(G_8) = 4$.

By Lemma 2.1 and 2.2, $\chi_\rho(G_n) = 4$. ■

Definition 2.8. [16] For a graph G , the splitting graph S' of G is obtained by adding a new vertex v' corresponding to each vertex v of G such that $N(v) = N(v')$.

Theorem 2.9. Let $K_{1,n}$, $n \geq 2$ be a star on n vertices. Then $\chi_\rho(S'(K_{1,n})) = 3$.

Proof: Let v_1, v_2, \dots, v_n be the pendent vertices and v be the apex vertex of $K_{1,n}$ and u, u_1, u_2, \dots, u_n be the added vertices corresponding to v, v_1, v_2, \dots, v_n to obtain $S'(K_{1,n})$.

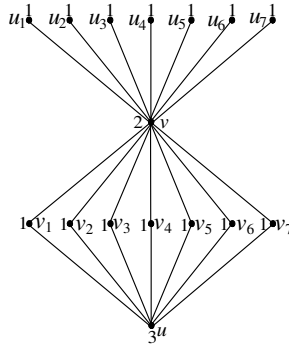


Figure 3: $\chi_\rho(S'(K_{1,7})) = 3$.

Color the vertices u_1, u_2, \dots, u_n and v_1, v_2, \dots, v_n with color 1. Assign color 2 to the vertex u and 3 to the vertex v . Thus $\chi_\rho(S'(K_{1,n})) \leq 3$. By Lemma 2.1 and 2.2, $\chi_\rho(S'(K_{1,n})) = 3$. ■

Definition 2.10. [12] The n -barbell graph is a simple graph obtained by connecting two copies of a complete graph K_n by a bridge and it is denoted by $B(K_n, K_n)$.

Theorem 2.11. For $B(K_n, K_n)$, $\chi_\rho(B(K_n, K_n)) = 2n - 2$.

Proof: An n -barbell graph contains two copies of K_n namely $1K_n$ and $2K_n$. Let uv be the connecting bridge where $u \in 1K_n$ and $v \in 2K_n$.

We give an algorithm to color $B(K_n, K_n)$ using exactly $2n - 2$ colors.

Procedure PACKING COLORING $B(K_n, K_n)$

Input: Barbell graph $B(K_n, K_n)$

Algorithm:

Step 1: Label the vertex u with color $2n - 2$.

Step 2: Label the vertex v with color $2n - 3$.

Step 3: Label any vertex of $1K_n$ and $2K_n$ with color 1.

Step 4: Label any vertex of $1K_n$ and $2K_n$ with color 2.

Step 5: Label the remaining vertices of $1K_n$ and $2K_n$ with color $3, 4, \dots, 2n - 4$.

Output: $\chi_\rho(B(K_n, K_n)) = 2n - 2$.

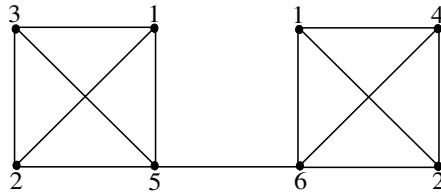


Figure 4: $\chi_\rho(B(K_4, K_4)) = 6$.

Proof of correctness: Let $G \simeq B(K_n, K_n)$. The diameter of G is 3, which implies $\rho_2(G) = 2$ and $\rho_i(G) \leq 1$ for all integer $i \geq 3$.

Case 1: If vertex u or v is labeled by color 1, then $n - 1$ vertices at distance greater than one from the vertex u or v form a subgraph, which is isomorphic to K_{n-1} . In this case, at most one vertex receives color 1, because $\rho_1(K_{n-1}) = 1$.

Case 2: If a vertex of G except u and v is labeled by color 1, then n vertices at distance more than one from that vertex received color 1 form a subgraph, which is isomorphic to K_n . In this case, at most one vertex receives color 1, because $\rho_1(K_n) = 1$.

So, $\rho_2(G) = 2$. There is a k -coloring of the graph G with $|X_1| = 2$ and $|X_2| = 2$. As for each of the other $(2n - 4)$ vertices, one new color is necessary, which implies that $\chi_\rho(B(K_n, K_n)) = 2n - 2$. ■

Definition 2.12. The graph obtained by joining a pendent edge at each vertex of a cycle C_n is called a crown graph denoted by $C_n \odot K_1$.

Theorem 2.13. For $C_n \odot K_1, n \geq 6$. $\chi_\rho(C_n \odot K_1) \leq 5$, when $n \equiv 0, 2 \pmod{4}$.

Proof: Let the vertices of C_n be u_1, u_2, \dots, u_n and the vertices of pendent edges be v_1, v_2, \dots, v_n where u_i and v_j are adjacent for all $i = j$.

Case 1: $n \equiv 0 \pmod{4}$. Color the vertices u_1, u_2, \dots, u_n with the sequence $1213, \dots$ starting from vertex u_1 in the clockwise sense and v_1, v_2, \dots, v_n with the sequence $4151, \dots$ starting from vertex v_1 in the clockwise sense. Thus $\chi_\rho(C_n \odot K_1) \leq 5$.

Case 2: $n \equiv 2 \pmod{4}$. Color the vertices u_1, u_2, \dots, u_n with the sequence $141315, \dots$ starting from vertex u_1 in the clockwise sense and v_1, v_2, \dots, v_n with the sequence $21, \dots$ starting from vertex v_1 in the clockwise sense. Thus $\chi_\rho(C_n \odot K_1) \leq 5$.

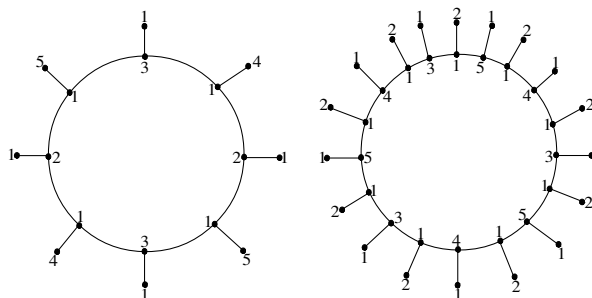


Figure 5: $\chi_\rho(C_8 \odot K_1) = 5$ and $\chi_\rho(C_{18} \odot K_1) = 5$.

■

Definition 2.14. The tadpole graph $T(m, n)$ is obtained from a cycle C_m and a path P_n , by joining one of the end vertices of P_n to a vertex of C_m .

Theorem 2.15. For $T(m, n)$, $n \geq 4$, $\chi_\rho(T(m, n)) = \begin{cases} 3, & \text{when } m \equiv 0 \pmod{4} \\ 4, & \text{when } m \not\equiv 0 \pmod{4} \end{cases}$.

Proof: Case 1: $m \equiv 0 \pmod{4}$.

Let u be a vertex in $T(m, n)$ with $\deg(u) = 3$. Color the vertices of cycle C_m with the sequence 2131, \dots starting from the vertex u in the clockwise sense. Let v be a vertex of path P_n which is adjacent to vertex u . Color the path P_n with the sequence 1213, \dots starting from the vertex v . Thus $\chi_\rho(T(m, n)) \leq 3$. By Lemma 2.1 and 2.2, $\chi_\rho(T(m, n)) = 3$.

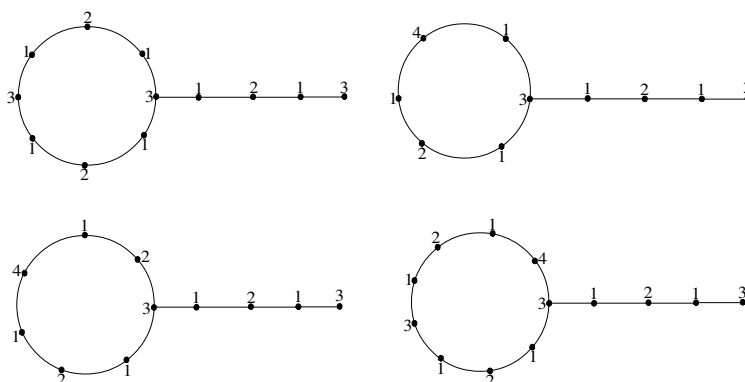


Figure 6: $\chi_\rho(T(8,5)) = 3, \chi_\rho(T(6,5)) = 4, \chi_\rho(T(7,5)) = 4$ and $\chi_\rho(T(9,5)) = 4$.

Case 2: $m \not\equiv 0 \pmod{4}$

Let u be a vertex in $T(m, n)$ with $\deg(u) = 3$. Color the vertices of cycle C_m with color sequence in the clockwise sense starting from vertex u with repeated blocks of 3,1,2,1 with an adjustment at the very end as shown below:

$$3121, 3121, \dots, 31214 \text{ when } n = 4r + 1$$

$$3121, 3121, \dots, 312141 \text{ when } n = 4r + 2$$

$$3121, 3121, \dots, 3121412 \text{ when } n = 4r + 3 \text{ starting from the vertex } u \text{ in the}$$

clockwise sense. Let v be the vertex of path P_n which is adjacent to vertex u . Color the path

P_n with the sequence $1213, \dots$ starting from the vertex v . Thus $\chi_\rho(T(m, n)) \leq 4$.

By Lemmas 2.1 and 2.2, $\chi_\rho(T(m, n)) = 4$. ■

Definition 2.16. [10] The lollipop graph $LP(n, k)$ is obtained from a complete graph K_k and a path P_{n-k+1} , by joining one of the end vertices of P_{n-k+1} to a vertex of K_k .

Theorem 2.17. For lollipop graph $LP(n, k)$, $\chi_\rho(LP(n, k)) = n$.

Proof: Let u be the vertex in $LP(n, k)$ with $\deg(u) = n$. Give color n to the vertex u and colors 1 to $n - 1$ to the remaining vertices in K_k .

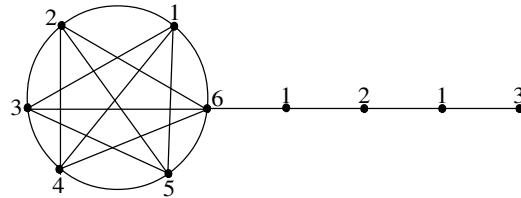


Figure 7: $\chi_\rho(LP(6, 2)) = 6$.

Let v be the vertex of path P_{n-k+1} which is adjacent to the vertex u . Color the path P_{n-k+1} with the sequence $1213, \dots$ starting from the vertex v . Thus $\chi_\rho(LP(n, k)) \leq n$.

By Lemma 2.1 and 2.3, $\chi_\rho(LP(n, k)) = n$. ■

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