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The superior complement in graphs

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Abstract

For distinct vertices u and v of a nontrivial connected graph G, we let $D_{u,v} = N[u] \cup N[v]$. We define a $D_{u,v}$ -walk as a u-v walk in G that contains every vertex of $D_{u,v}$. The superior distance $d_D(u,v)$ from u to v is the length of a shortest $D_{u,v}$ -walk. For each vertex $u \in V(G)$, define $d_D^-(u) = \min\{d_D(u,v) : v \in V(G) - \{u\}\}$. A vertex $v(\neq u)$ is called a *superior neighbor* of u if $d_D(u,v) = d_D^-(u)$. In this paper we define the concept of superior complement of a graph G as follows: The superior complement of a graph G is denoted by \overline{G}_D whose vertex set is as in G. For a vertex u, let $A_u = \{v \in V(G) : d_D(u,v) \ge d_D^-(u) + 1\}$. Then u is adjacent to all the vertices $v \in A_u$ in \overline{G}_D . The main focus of this paper is to prove that there is no relationship between the superior complement \overline{G}_D of G.

Key words: Superior distance, superior radius, superior diameter, superior neighbor, superior complement.

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1 Introduction

By a graph we mean a nontrivial undirected graph without loops and multiple edges. As usual V(G) denotes the set of vertices of a graph G and E(G) denotes the set of edges of G. The distance between the vertices u and v is the length of a shortest u-v path in G. The distance to a vertex farthest from a vertex u is the eccentricity e(u) of the vertex u in G. The minimum among the eccentricities is called the radius r(G) of G and the maximum among the eccentricities is called the diameter d(G) of G. By a neighbor of a vertex u, we mean, any vertex whose distance from u is minimum. If v is a vertex distinct from u whose distance from uis minimum, then this distance d(u, v) must be 1. For each vertex $u \in V(G)$, define $d^{-}(u) = min\{d(u, v) : v \in V(G) - \{u\}\}$. A vertex $v \neq u$ is called a neighbor of u if $d^{-}(u) = d(u, v)$. Since $d^{-}(u) = 1$ for all $u \in V(G)$, this is equivalent to the standard definition of neighbor. For graph theoretic notation and terminology, we follow [1].

If X and Y are two cities, then for a taxi driver, the distance between X and Y is the actual distance between the two cities. However for a driver of a passenger bus, the distance between the same two cities is greater than the usual distance since he has to visit important places in and around the two cities to pickup and drop off passengers. So a driver of a passenger bus has to find a shortest route that begins at X, ends at Y, and passes through each of the neighboring places of X and Y.

Kathiresan and Marimuthu [2] defined a variation of distance that models the bus route just described. For two vertices u and v of G, let $D_{u,v} = N[u] \cup N[v]$. They define a $D_{u,v}$ -walk as a u-v walk in G that contains every vertex of $D_{u,v}$. The superior distance $d_D(u, v)$ from u to v is the length of a shortest $D_{u,v}$ -walk. If u and v are in the different components of a disconnected graph, then we define $d_D(u, v) = \infty$.

The superior eccentricity $e_D(u)$ of a vertex u is the superior distance to a vertex farthest from u. A vertex v is called a superior eccentric vertex of a vertex u if $e_D(u) = d_D(u, v)$. The minimum among the superior eccentricities in a graph Gis called the superior radius $r_D(G)$ of G and the maximum among the superior eccentricities in a graph G is called the superior diameter $d_D(G)$ of G.

One variation of domination is seen in [3] where the concept of superior domination in graphs is defined. In the standard definition of domination in a graph, a vertex u dominates itself and each of its neighbors. For each vertex $u \in V(G)$, define $d_D^-(u) = \min\{d_D(u, v) : v \in V(G) - \{u\}\}$. A vertex $v(\neq u)$ is called a superior neighbor of u if $d_D(u, v) = d_D^-(u)$. The superior neighborhood $N_D(u)$ of a vertex u is the set of all superior neighbors of u; and its closed superior neighborhood is $N_D[u] = N_D(u) \cup \{u\}$. A vertex u is said to superior dominate a vertex v if v is a superior neighbor of u. A set S of vertices of G is called a superior dominating set if every vertex of V - S is superior dominated by some vertex of S. The complement \overline{G} of a graph G is the graph with vertex set as in G and two vertices u and v are adjacent in \overline{G} if and only if $d(u, v) \geq 2$. As there is a complement of a graph G with respect to usual distance, it is natural to define a new complement of a graph with respect to the superior distance. There are some cases in which the usual complement and the superior complement coincide. For example $\overline{C_4} = \overline{(C_4)}_D$.

In this paper, we define the concept of superior complement in graphs as follows: The superior complement of a graph G is denoted by \overline{G}_D whose vertex set is as in G. For each vertex $u \in V(G)$, let $A_u = \{v \in V(G) : d_D(u, v) \ge d_D^-(u) + 1\}$. Then u is adjacent to all the vertices v in A_u in \overline{G}_D .

Next we provide one result which is found in [4].

Proposition A.[4] Let G be a nontrivial connected graph with $N[u] \cup N[v] = V(G)$ for any two distinct vertices u and v. Then $d_D(u, v) = n - 1$ for any two distinct vertices u and v if and only if G is hamiltonian connected.

2 The Superior Complement

Observation 2.1. For any connected graph G of order n, $d_D(G) \le 2n - 3$.

The sharpness of the upper bound is satisfied by the graphs stars and the double stars.

Theorem 2.2. Let G be a graph of order n. Then $\overline{G}_D = K_n$ if and only if G is totally disconnected.

Proof. Assume that $\overline{G}_D = K_n$. We claim that $G = \overline{K}_n$. Suppose G has at least one edge uv (say). Then by the definition of \overline{G}_D , $uv \notin \overline{G}_D$, a contradiction.

If $G = \overline{K}_n$, then $\overline{G}_D = K_n$ is followed from the definition.

Proposition 2.3. If $G = K_{1,n-1}$, then $\overline{G}_D = G$.

Observation 2.4. Let G be a graph of order n. Then $\overline{G}_D = \overline{G}$ if and only if $N(u) = N_D(u)$ for all $u \in V(G)$.

Theorem 2.5. Let G be a nontrivial connected graph of order n such that $N[u] \cup N[v] = V(G)$ for each pair of distinct vertices u and v. Then G is hamiltonian connected if and only if $\overline{G}_D = \overline{K_n}$.

Proof. Assume that G is hamiltonian connected. Then by Proposition A, $d_D(u, v) = n - 1$ for each pair of distinct vertices u and v. Thus for a vertex u, every vertex except u is a superior neighbor of u and hence $\overline{G}_D = \overline{K}_n$.

Conversely assume that $\overline{G}_D = \overline{K}_n$. It is enough if we prove that $d_D(u, v) = n - 1$ for each pair of distinct vertices u and v. Since $\overline{G}_D = \overline{K}_n$, for each vertex u, all the vertices except u are the superior neighbor of u in G. Thus $d_D(u, v) = l$ for all $u, v \in V(G)$ in G. We claim that n - 1 = l.

Suppose there exists a pair of vertices u and v such that $d_D(u, v) = l < n - 1$. Then v is a superior neighbor of u and there are some vertices in A_u . Thus there is at least one edge in \overline{G}_D , a contradiction.

Suppose there exist a pair of vertices u and v such that $d_D(u, v) = l > n - 1$. Then u and v are non superior neighbor to each other. Thus $uv \in E(\overline{G}_D)$. This is a contradiction.

Lemma 2.6. Let G be a hamiltonian connected graph of order $n \ge 2$ such that $N[u] \cup N[v] = V(G)$ for each pair of distinct vertices u and v. Then $d_D(\overline{G}_D) = \infty$.

Proof. Let G be a Hamiltonian connected graph of order $n \ge 2$ such that $N[u] \cup N[v] = V(G)$ for each pair of vertices u and v. Then by Proposition A, $d_D(u, v) = n - 1$, for each pair of vertices u and v. Then $d_D(G) = r_D(G) = n - 1$.

It follows that $d_D^-(u) = n - 1$ for all $u \in V(G)$. Thus for each vertex u, every vertex of G other than u is a superior neighbor of u. Thus $\overline{G}_D = \overline{K}_n$ and hence $d_D(\overline{G}_D) = \infty$.

Proposition 2.7. For every $n \ge 5$, there exists a connected graph G of order n such that $\overline{G}_D = \overline{G}$.

Proof. Let G be any cycle $C_n, n \ge 5$ on n vertices. Let u be any vertex of C_n .

Then $d_D(u, v) = 5$ if v is adjacent to u and $d_D(u, v) > 5$ if v is nonadjacent to u in C_n .

Thus $N_D(u) = N(u)$ for all $u \in V(C_n)$. By Observation 2.4, $\overline{G}_D = \overline{G}$.

Remark 2.8. If $G = K_n$, then $\overline{G}_D = \overline{G}$ for any n.

Remark 2.9. Also \overline{G}_D is a hamiltonian connected graph such that $N[u] \cup N[v] = V(\overline{G}_D)$ for each pair of vertices u and v. By Proposition A, $d_D(u, v) = n - 1$ for any u and v.

Corollary 2.10. There exists a graph G such that $\overline{H}_D \neq G$ where $H = \overline{G}_D$.

The following results are found in [1].

Theorem 2.11. If G is a simple graph with diameter at least 3, then \overline{G} has diameter at most 3.

Theorem 2.12. If G is a simple graph with diameter at least 4, then \overline{G} has diameter at most 2.

Even though there is a relation between the diameter of G and the diameter of \overline{G} in the usual distance, we show that there is no such relationship between the superior diameter $d_D(G)$ of a graph G and the superior diameter of the superior complement \overline{G}_D of G.

Lemma 2.13. For each positive integer $n \ge 4$, there exists a connected graph G of order n such that $d_D(\overline{G}_D) = 2n - 4$ where $d_D(G) = n$.

Proof. Consider the graph K_{n-1} for any $n \ge 4$. Let u be the new vertex and let $u_i, i = 1, 2, ..., n - 1$ be the vertices of K_{n-1} . Now join u with any one of the vertices of K_{n-1} . Assume that u is adjacent to u_1 . Let the resulting graph be G. Then $d_D(u, u_1) = n$ in G. The superior eccentric vertex of each vertex of G except u_1 is u_1 . Thus $e_D(u) = n$ for all u in G. $d_D(u, u_i) = n - 1$ for all i = 2, 3, ..., n - 1; $d_D(u_i, u_j) = n - 2$ for all $i \ne j, i, j = 2, 3, ..., n - 1$. Thus each $u_i, i = 2, 3, ..., n - 1$ is adjacent to u and u_1 in \overline{G}_D . Also u is adjacent to u_1 in \overline{G}_D . Now, $d_D(u, u_1) = 2n - 4$ and $d_D(u_1, u) = 2n - 4$ in \overline{G}_D . $d_D(u_i, u) = 2n - 5$ for all $u_i, i = 2, 3, ..., n - 1$; $d_D(u_i, u_j) = 4$ for all $i \neq j, i, j = 2, 3, ..., n - 1$ in \overline{G}_D . Thus $d_D(\overline{G}_D) = 2n - 4$.



Figure 1: The graphs G and \overline{G}_D with superior eccentricity with five vertices

Remark 2.14. The direct verification shows that the result is not true for $n \leq 3$.

Lemma 2.15. Let G be a complete bipartite graph $K_{m,n}$. Then the followings hold.

(a)
$$\overline{G}_D = \begin{cases} K_m \cup K_n & \text{if } m = n \\ K_m + \overline{K_n} & \text{if } m < n \end{cases}$$

(b)
$$d_D(G) = d_D(\overline{G}_D) = 2n \text{ if } n > m = 2.$$

Proof. (a) Let V_1 and V_2 be the bipartition of the vertex set of G such that $|V_1| = m$ and $|V_2| = n$. Assume that m = n. Let u and v be the vertices in the same partite set. Then $d_D(u, v) = 2m$. If u and v belong to different sets, then $d_D(u, v) = 2m - 1$. Thus for each vertex $u \in V_1$, every vertex in V_2 is a superior neighbor of u. Thus uis adjacent to all other vertices in V_1 in \overline{G}_D . Similarly each vertex u in V_2 is adjacent to all other vertices in V_2 in \overline{G}_D . Hence $\overline{G}_D = K_m \cup K_n$, a disconnected graph.

Now, assume that m < n. Then every vertex of V_1 is a superior eccentric vertex of all other vertices in V_1 . Let $u \in V_1$ and $v \in V_2$. Then $d_D(u,v) = 2n - 1$; $d_D(u,v) = 2n$ if $u, v \in V_1$; $d_D(u,v) = 2m$ if $u, v \in V_2$ in G. Thus $\overline{G}_D = K_m + \overline{K_n}$.

(b) Assume that n > m = 2. Then $d_D(u, v) = 2n$ if $u, v \in V_1$; $d_D(u, v) = 2n - 1$ if $u \in V_1$ and $v \in V_2$; $d_D(u, v) = 2m$ if $u, v \in V_2$ in G. This is also true in \overline{G}_D , since $\overline{G}_D = K_2 + \overline{K_n}$, a connected graph.

Lemma 2.16. For each odd l, there exists a connected graph G such that $d_D(G) = d_D(\overline{G}_D) = l$.

Proof. Consider the graph $K_{m,n}$. Assume that m = 1 and $n \ge 1$. When n = 1 the result is trivial. For otherwise, let u be the vertex of degree n in G. Then $d_D(u, v) = 2n-1$ for all $v \in V_2$; $d_D(w, v) = 2$ if $w, v \in V_2$. Also $d_D(v, u) = 2n-1$ for all $v \in V_2$. Clearly, $G = \overline{G}_D$. Thus $d_D(G) = d_D(\overline{G}_D) = 2n-1 = l$ (say).

Proposition 2.17. There exist distinct graphs G and H with $\overline{G}_D \cong \overline{H}_D$.

Proof. Let G be the graph constructed in Lemma 2.13. Take $H = K_{2,n-2}$. Then $\overline{G}_D = \overline{H}_D$.

Lemma 2.18. There exist graphs G such that $d_D(\overline{G}_D) < d_D(G)$.

Proof. Let C_m be a cycle, $m = 3, 4, 5, \ldots, 8$. Attach a path of length 2 at any vertex of the cycle C_m and let the resulting graph be G_m . It is easy to verify that $d_D(G_m) \le 10; d_D(\overline{G_m})_D \le 9$ and $d_D(\overline{G_m})_D < d_D(G)$ for all G_m .

Theorem 2.19. There is no relationship between the superior diameter $d_D(G)$ of a graph G and the superior diameter of the superior complement \overline{G}_D of G.

Proof. Follows from Lemmas 2.13, 2.15, 2.16 and 2.17.

Open Problems:

- 1. Characterize all graphs G for which $\overline{G}_D = G$.
- 2. Characterize all graphs G for which $\overline{H}_D = G$ where $H = \overline{G}_D$.
- 3. Given any natural numbers a and b (large enough), does there exist a graph G such that $d_D(G) = a$ and $d_D(\overline{G}_D) = b$?

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