# A New Inexact Inverse Subspace Iteration for Generalized Eigenvalue Problems 

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#### Abstract

In this paper, we represent an inexact inverse subspace iteration method for computing a few eigenpairs of the generalized eigenvalue problem $A x=\lambda B x[\mathrm{Q}$. Ye and P . Zhang, Inexact inverse subspace iteration for generalized eigenvalue problems, Linear Algebra and its Application, 434 (2011) 1697-1715 ]. In particular, the linear convergence property of the inverse subspace iteration is preserved.


Keywords: Eigenvalue problem; inexact inverse iteration; subspace iteration; inner-outer iteration; approximation; convergence.

## 1. Introduction

We want to compute a few eigenpairs of the generalized eigenvalue problem $A x=$ $\lambda B x$. (1) The eigenvalues sought may be those in the extreme part of the spectrum or in the interior of the spectrum near certain given point. These types of problems arise in many engineering and scientific applications. In such applications, the matrices involved are often large and sparse. A brief description of some developed method is given in
[2] . If the extreme eigenvalues are not well-separated or if the eigenvalues sought are in the interior of the spectrum, a shift-and invert (or inverse) transformation is combined with one of the eigenproblem solvers to speed up the convergence. The use of the shift-and-invert transformation requires solving a linear system at each iterative step. Solving the linear systems by a direct method such as the QR factorization can be impractical or expensive if the dimension of the matrices is large or the matrices are not explicitly available. Alternatively, iterative method can be employed, which will be called the inner iteration while the original iterative algorithm will be called the outer iteration. The inner iteration produces, a situation that may arise in other applications as well. The question then is how the accuracy in the inner iterations affects the convergence behavior (or convergence speed) of the outer iteration as compared with the exact case. Related to this, a challenging problem in implementations is how to efficiently choose an appropriate stopping threshold for the inner iteration so that the convergence characteristic of the outer iteration can be preserved. This is a problem that has been discussed for several methods, such as inexact Krylov subspace method [6, 10, 17] inexact

[^0]inverse iteration, $[1,4,8,12,13]$, the rational Arnoldi algorithm [9, 11, 16], inexact Rayleigh Quotient-Type methods [3, 7, 14] and the Jacobi-Davidson method [15]. While these works demonstrate that the innerouter iteration technique may be an effective way for implementing some of these methods for solving the large scale eigenvalue problem, the more efficient Krylov subspace projection methods tend to require quite accurate matrix-vector products to preserve their convergence characteristic.

## 2. Inexact inverse subspace iteration

We consider computing the $p$ smallest eigenvalues of the generalized eigenvalue problem $A x=\lambda B x$ by applying the standard subspace iteration to $A^{-1} B$, called inverse subspace iteration. we state the standard algorithm as follows.

Algorithm 1. Inverse subspace iteration for $A x=\lambda B x$
(1) Input: $X_{0} \in C^{n \times p}$ with $X_{0}^{*} X_{0}=I$;
(2) For $k=0,1, \cdots$ until convergence
(3) $Y_{k+1}=A^{-1} B X_{k}$;
(4) $Y_{k+1}=X_{k+1} R_{k+1}$ (QR-factorization).
(5) End.

We try to use this algorithm when a problems where direct solution of $A^{-1}$ is impractical or inefficient because $A$ is too large or is not explicitly available. In these cases, an iterative method can be used to solve the linear systems $A Y_{k+1}=B X_{k}$, called the inner iterations, while the subspace iteration itself is called the outer iteration. At each step of outer iteration, to solve $Y_{k+1}$ for $A Y_{k+1}=B X_{k}$, the previous iterate $Y_{k}$ can be used as an initial approximation. Then we solve

$$
\begin{equation*}
A D_{k}=B X_{k}-A Y_{k} \tag{1}
\end{equation*}
$$

approximately, i.e. in this case we find $D_{k}$ such that

$$
\begin{equation*}
\left\|E_{k}\right\|_{2}=\left\|\left(B X_{k}-A Y_{k}\right)-A D_{k}\right\|_{2}<\epsilon_{k} \tag{2}
\end{equation*}
$$

where $E_{k}:=\left(B X_{k}-A Y_{k}\right)-A D_{k}$ and $k$ is some given threshold. From $D_{k}$, then $Y_{k+1}=Y_{k}+D_{k}$.

Now we try to analyze the convergence characteristic of the subspace iteration under inexact solves.

Obviously, the amount of work required to solve (2) is proportional to $\| B X_{k}-$ $A Y_{k} \| / \epsilon_{k}$. Our analysis later leads to the use of linearly decreasing $k$, let $\epsilon_{k}=a r^{k}$ for some positive a and $r<1$. However, $\epsilon_{k}$ is decreasing, the amount of work does not increase as $B X_{k}-A Y_{k}$ will be decreasing at the same rate as well. Thus, the stopping threshold required for inner iterations is effectively a constant.

Algorithm 2. Inexact inverse subspace iteration for $\mathrm{Ax}=\mathrm{Bx}$
(1) Input: $X_{0} \in C^{n \times p}$ with $X_{0}^{*} X_{0}=I$; threshold parameter $\epsilon_{k}$; set $Y_{0}=0$;
(2) For $k=0,1, \ldots$ until convergence
(3) $Z_{k}=B X_{k}-A Y_{k}$;
(4) Solve $A D_{k}=Z_{k}$ such that $E_{k}=Z_{k}-A D_{k}$ satisfies (3);
(5) $Y_{k+1}=Y_{k}+D_{k}$;
(6) $Y_{k+1}=\bar{X}_{k+1} \bar{R}_{k+1}$ (QR-factorization);
(7) For $j=1, \ldots, p$
(8) $\left[y_{\max }, i_{\max }\right]=\max \left(\left|\bar{X}_{k+1}(:, j)\right|\right)$;
(9) $X_{k+1}(:, j)=\operatorname{sign}\left(\bar{X}_{k+1}\left(i_{\max }, j\right)\right) * X_{k+1}(:, j)$;
(10) $R_{k+1}(j,:)=\operatorname{sign}\left(\bar{X}_{k+1}\left(i_{\max }, j\right)\right) * R_{k+1}(j,:)$.
(11) End.
(12) End.

For more details please refer to [4]

## 3. Convergence analysis

We discuss convergence of the subspace spanned by $X_{k}$ for the inexact inverse subspace algorithm. Let $\lambda_{1}, \lambda_{2}, \cdots, \lambda_{n}$ be the eigenvalues of $B^{-1} A$ ordered such that

$$
0<\left|\lambda_{1}\right| \leqslant \cdots \leqslant\left|\lambda_{p}\right|<\left|\lambda_{p+1}\right| \leqslant \cdots \leqslant\left|\lambda_{q}\right|<\left|\lambda_{q+1}\right| \leqslant \cdots\left|\lambda_{n}\right|
$$

and $v_{1}, \ldots, v_{n}$ be the corresponding eigenvectors. Suppose that we want to compute the $p$ smallest eigenpairs in absolute value, i.e., $\lambda_{1}, \lambda_{2}, \cdots, \lambda_{p}, \lambda_{p+1}, \cdots, \lambda_{q}$. We are following eigenvalues of $p+1$ to $q$ (with assumption ordered ) so we have made the assumptions that $\rho:=\left|\lambda_{p}\right| /\left|\lambda_{p+1}\right|<1, \dot{\rho}:=\left|\lambda_{q}\right| /\left|\lambda_{q+1}\right|<1$.

Throughout this work, we assume that $B^{-1} A$ is diagonalizable. Let $V=$ $\left[v_{1}, \ldots, v_{n}\right], U=(B V)^{-H}$,then

$$
\begin{aligned}
& U^{H} A=\Lambda U^{H} B, A V=B V \Lambda \text { where } \Lambda=\left(\begin{array}{ccc}
\Lambda_{1} & 0 & 0 \\
0 & \Lambda_{2} & 0 \\
0 & 0 & \Lambda_{3}
\end{array}\right) \\
& \Lambda_{1}=\left(\begin{array}{ccc}
\lambda_{1} & \cdots & 0 \\
\vdots & \ddots & \vdots \\
0 & \cdots & \lambda_{p}
\end{array}\right), \Lambda_{2}=\left(\begin{array}{ccc}
\lambda_{p+1} & \cdots & 0 \\
\vdots & \ddots & \vdots \\
0 & \cdots & \lambda_{q}
\end{array}\right), \Lambda_{1}=\left(\begin{array}{ccc}
\lambda_{q+1} & \cdots & 0 \\
\vdots & \ddots & \vdots \\
0 & \cdots & \lambda_{n}
\end{array}\right)
\end{aligned}
$$

let $U=\left(U_{1}, U_{2}, U_{3}\right), V=\left(V_{1}, V_{2}, V_{3}\right)$, where $U_{1} \in^{n \times p}, U_{2} \in^{n \times(p+q)}, U_{3} \in$ $n \times\left(n-(p+q), V_{1} \in^{n \times p}, V_{2} \in^{n \times(p+q)} V_{3} \in^{n \times(n-(p+q)}\right.$, then $U_{i}^{H} A=\Lambda_{i} U_{i}^{H} B, A V_{i}=$ $B V_{i} \Lambda_{i}$, Consider Algorithm 2 now. Define $X_{k}^{(i)}=U_{i}^{H} B X_{k}$. Since $U_{i}^{H} B V_{j}=\delta_{i j} I$ and $U_{i}^{H} A V_{j}=\delta_{i j} \Lambda_{i}$ where $\delta_{i j}$ is the Kronecker symbol, then

$$
X_{k}=\sum_{i=1}^{3} V_{i} X_{k}^{(i)}
$$

If $X_{k}^{(1)}$ is invertible, we define $t_{k}:=\left\|X_{k}^{(2)}\left(X_{k}^{(1)}\right)^{-1}\right\|_{2}$ and $t_{k^{\prime}}:=\left\|X_{k}^{(3)}\left(X_{k}^{(1)}\right)^{-1}\right\|_{2}$ Clearly, $t_{k}$ and $t_{k^{\prime}}$ is a measures of the approximation of the column space of $X_{k}$ to the column space of $V_{1}$. Indeed, the following proposition relates $t_{k}$ and $t_{k^{\prime}}$ to other measures of subspace approximation.
Proposition 3.1 Assume that $X_{k}^{(1)}$ is invertible and $t_{k}$ and $t_{k^{\prime}}$ is defined as above. Then

$$
\begin{equation*}
\frac{\binom{t_{k}}{t_{k^{\prime}}}}{\left\|V^{-1}\right\|_{2}} \leqslant\left\|X_{k}\left(X_{k}^{(1)}\right)^{-1}-V_{1}\right\|_{2} \leqslant\|V\|_{2}\left(t_{k}+t_{k^{\prime}}\right) \tag{3}
\end{equation*}
$$

and

$$
\sin \angle\left(\chi_{k}, \nu_{1}\right) \leqslant\|V\|_{2}\left\|R^{-1}\right\|_{2}\left(t_{k}+t_{k^{\prime}}\right)
$$

where $\angle\left(X_{k}, V_{1}\right)$ is the largest canonical angle between $X_{k}=R\left(X_{k}\right)$ and $V_{1}=$ $R\left(V_{1}\right)$, and $R$ is defined by the $Q R$-factorization $V_{1}=W R$ of $V_{1}$.

Proof From $X_{k}=V_{1} X_{k}^{(1)}+V_{2} X_{k}^{(2)}+V_{3} X^{(3)}$, we have

$$
\begin{aligned}
& X_{k}\left(X_{k}^{(1)}\right)^{-1}=V_{1}+V_{2} X_{k}^{(2)}\left(X_{k}^{(1)}\right)^{-1}+V_{2} X_{k}^{(3)}\left(X_{k}^{(1)}\right)^{-1} \\
& \left\|X_{k}\left(X_{k}^{(1)}\right)^{-1}-V_{1}\right\|_{2}=\left\|V_{2} X_{k}^{(2)}\left(X_{k}^{(1)}\right)^{-1}+V_{3} X_{k}^{(3)}\left(X_{k}^{(1)}\right)^{-1}\right\|_{2} \\
& \leqslant\left\|V_{2}\right\|_{2}\left\|X_{k}^{(2)}\left(X_{k}^{(1)}\right)^{-1}+\right\| V_{3}\left\|_{2}\right\| X_{k}^{(3)}\left(X_{k}^{(1)}\right)^{-1}\left\|_{2} \leqslant\right\| V_{2} \|_{2}\left(t_{k}+t_{k^{\prime}}\right)
\end{aligned}
$$

and

$$
\left\|V_{2} X_{k}^{(2)}\left(X_{k}^{(1)}\right)^{-1}+V_{3} X_{k}^{(3)}\left(X_{k}^{(1)}\right)^{-1}\right\|_{2}=\left\|V\binom{X_{k}^{(2)}\left(X_{k}^{(1)}\right)^{-1}}{X_{k}^{(3)}\left(X_{k}^{(1)}\right)^{-1}}\right\|_{2} \geqslant \frac{\binom{t_{k}}{t_{k^{\prime}}}}{\left\|V^{-1}\right\|_{2}}
$$

(4) is proved.

Let $X_{k}^{\perp}$ be such that $\left(X_{k}, X_{k}^{\perp}\right)$ is an $n \times n$ orthogonal matrix. Then the sine of the largest canonical angle between $\chi_{k}=R\left(X_{k}\right)$ and $\nu_{1}=R\left(V_{1}\right)$ is (see [5] for the definition)

$$
\begin{aligned}
\sin \angle\left(\chi_{k}, \nu_{1}\right) & =\left\|\left(X_{k}^{\perp}\right)^{H} w\right\|_{2} \\
& =\left\|\left(X_{k}^{\perp}\right)^{H}\left(w-X_{k}\left(X_{k}^{(1)}\right)^{-1} R^{-1}\right)\right\|_{2} \\
& =\left\|\left(X_{k}^{\perp}\right)^{H}\left(V_{2} X_{k}^{(2)}\left(X_{k}^{(1)}\right)^{-1} R^{-1}+V_{3} X_{k}^{(3)}\left(X_{k}^{(1)}\right)^{-1} R^{-1}\right)\right\|_{2} \\
& \leqslant\|V\|_{2}\left\|R^{-1}\right\|_{2}\left(t_{k}+t_{k^{\prime}}\right)
\end{aligned}
$$

It is clear that $t_{k}$ and $t_{k^{\prime}}$ are measures of the approximation of the column space of $X_{k}$. We shall next discuss the convergence of $t_{k}$ and $t_{k^{\prime}}$.

Lemma 3.2 For Algorithm $2\left\|E_{k}\right\|_{2}<\left\|B^{-1}\right\|_{2}^{-1}$ then $Y_{k+1}$ has full column rank.
Proof From the algorithm, we have $A Y_{k+1}=B X_{k}+E_{k}$. Therefore

$$
X_{k}^{H} B^{-1} A Y_{k+1}=I+X_{k}^{H} B^{-1} E_{k}
$$

Since

$$
\begin{gathered}
\left\|X_{k}^{H}\right\|_{2}=1 \\
\left\|X_{k}^{H} B^{-1} E_{k}\right\|_{2} \leqslant\left\|B^{-1}\right\|_{2}\left\|E_{k}\right\|_{2}<1
\end{gathered}
$$

$X_{k}^{H} B^{-1} A Y_{k+1}$ is invertible. Thus $Y_{k+1}$ has full column rank.
From now on, we shall assume that $\epsilon_{k} \leqslant\left\|B^{-1}\right\|_{2}^{-1}$, so that all $Y_{k}$ will have full column rank and Algorithm 2 will be well defined.

Lemma 3.3 For Algorithm 2, if $X_{k}$ is invertible, then

$$
\left\|\left(X_{k}^{(1)}\right)^{-1}\right\|_{2} \leqslant\|V\|_{2}\left(1+t_{k}+t_{k^{\prime}}\right)
$$

Proof Proof. Since $X_{k}$ has orthonormal columns, we have

$$
X_{k}=V_{1} X_{k}^{(1)}+V_{2} X_{k}^{(2)}+V_{3} X_{k}^{(3)}
$$

from which it follows that $\left(X_{k}^{(1)}\right)^{-1}$

$$
X_{k}\left(X_{k}^{(1)}\right)^{-1}=V_{1}+V_{2} X_{k}^{(2)}\left(X_{k}^{(1)}\right)^{-1}+V_{3} X_{k}^{(3)}\left(X_{k}^{(1)}\right)^{-1}
$$

since $X_{k}$ have orthonormal columns, we have from property of orthonormal matrix (if $Q$ is a orthonormal matrix then for every $x$ we have $\|Q x\|_{2}=\|x\|_{2}$ ) let

$$
\left\|\left(X_{k}^{(1)}\right)^{-1}\right\|_{2}=\left\|X_{k}\left(X_{k}^{(1)}\right)^{-1}\right\|_{2}
$$

we have the following
$\left\|\left(X_{k}^{(1)}\right)^{-1}\right\|_{2}=\left\|V_{1}+V_{2} X_{k}\left(X_{k}^{(1)}\right)^{-1}+V_{3} X_{k}^{(3)}\left(X_{k}^{(1)}\right)^{-1}\right\|_{2} \leqslant\left\|V_{1}\right\|_{2}+t_{k}\left\|V_{2}\right\|_{2}+t_{k^{\prime}}\left\|V_{3}\right\|_{2}$

$$
\leqslant\left(1+t_{k}+t_{k^{\prime}}\right)\|V\|_{2}
$$

## Lemma 3.4

Let rho $=\left|\lambda_{p}\right| /\left|\lambda_{p+1}\right|<1$ and assume that $X_{k}^{(1)}$ and $X_{k+1}^{(1)}$ are nonsingular. If $\|V\|_{2}\|U\|_{2}\left(1+t_{k}+t_{k^{\prime}}\right) \epsilon_{k}<1$, then

$$
t_{k+1} \leqslant \rho\left(t_{k}+t_{k^{\prime}}\right)+\frac{\rho\|V\|_{2}\|U\|_{2}\left(1+t_{k}+t_{k^{\prime}}\right)^{2} \epsilon_{k}}{1-\|V\|_{2}\|U\|_{2}\left(1+t_{k}+t_{k^{\prime}}\right) \epsilon_{k}}
$$

Proof From the algorithm we know that $A Y_{k+1}=B X_{k}+E_{k}$ and $Y_{k+1}=$ $X_{k+1} R_{k+1}$. Since $Y_{k+1}$ has full column rank, $R_{k+1}$ is invertible. Then

$$
A X_{k+1} R_{k+1} R_{k+1}^{(-1)}=B X_{k} R_{k+1}^{(-1)}+E_{k} R_{k+1}^{(-1)}
$$

Multiplying $U_{i}^{H}$ on the equation above, we have

$$
\begin{gathered}
\left\{\begin{array}{c}
X_{k}^{(i)}=U_{i}^{H} B X_{k} \\
U_{i}^{H} A=\Delta_{i} U_{i}^{H} B
\end{array}\right. \\
U_{i}^{H} A X_{k+1}=U_{i}^{H} B X_{k} R_{k+1}^{(-1)}+U_{i}^{H} E_{k} R_{k+1}^{(-1)} \\
\Lambda_{i} U_{i}^{H} B X_{k+1}=X_{k}^{(i)} R_{k+1}^{(-1)}+U_{i}^{H} E_{i} R_{k+1}^{(-1)}
\end{gathered}
$$

where

$$
\begin{gather*}
U_{i}^{H} B X_{k+1}=\Lambda_{i}^{(-1)} X_{k}^{(i)} R_{k+1}^{(-1)}+\Lambda_{i}^{(-1)} U_{i}^{H} E_{k} R_{k+1}^{(-1)} \\
X_{k+1}^{(i)}=\Lambda_{i}^{(-1)} X_{k}^{(i)} R_{k+1}^{(-1)}+\Lambda_{i}^{(-1)} \Delta_{k}^{(i)} \tag{4}
\end{gather*}
$$

$\Delta_{k}^{(i)}=U_{i}^{H} E_{k} R_{k+1}^{(-1)}$. let $i=2$

$$
\begin{aligned}
X_{k+1}^{(2)} & =\Lambda_{2}^{(-1)} X_{k}^{(2)} R_{k+1}^{(-1)}+\Lambda_{2}^{(-1)} \Delta_{k}^{(2)} \\
X_{k+1}^{(2)}\left(X_{k+1}^{(1)}\right)^{(-1)} & =\left(\Lambda_{2}^{(-1)} X_{k}^{(2)} R_{k+1}^{(-1)}+\Lambda_{2}^{(-1)} \Delta_{k}^{(2)}\right)\left(X_{k+1}\right)^{(-1)} \\
& =\left(\Lambda_{2}^{(-1)} X_{k}^{(2)} R_{k+1}^{(-1)}\right)\left(X_{k+1}\right)^{(-1)}+\Lambda_{2}^{(-1)} \Delta_{k}^{(2)}\left(X_{k+1}\right)^{(-1)} \\
& \left.=\Lambda_{2}^{(-1)} X_{k}^{(2)}\left(X_{k}^{(1)}\right)^{(-1)} \Lambda_{1}\right)\left(\left(\left(X_{k}\right)^{(1)}\right) \Lambda_{1}\right)^{-1} R_{k+1}^{(-1)}\left(X_{k+1}\right)^{(-1)} \\
& +\Lambda_{2}^{(-1)} \Delta_{k}^{(2)}\left(X_{k+1}\right)^{(-1)} \\
& \left.=\Lambda_{2}^{(-1)} X_{k}^{(2)}\left(X_{k}^{(1)}\right)^{(-1)} \Lambda_{1}\right) \Lambda_{1}^{-1}\left(\left(X_{k}\right)^{(1)}\right)^{-1} R_{k+1}^{(-1)}\left(X_{k+1}\right)^{(-1)} \\
& +\Lambda_{2}^{(-1)} \Delta_{k}^{(2)}\left(X_{k+1}\right)^{(-1)}
\end{aligned}
$$

and now in the (4) let $i=1$

$$
\begin{equation*}
X_{k+1}^{(1)}=\Lambda_{1}^{(-1)} X_{k}^{(1)} R_{k+1}^{(-1)}+\Lambda_{1}^{(-1)} \Delta_{k}^{(1)} \tag{5}
\end{equation*}
$$

$$
\Lambda_{1}^{(-1)} X_{k}^{(1)} R_{k+1}^{(-1)}=X_{k+1}^{(1)}-\Lambda_{1}^{(-1)} \Delta_{k}^{(1)}
$$

$$
\begin{aligned}
X_{k+1}^{(2)}\left(X_{k+1}^{(1)}\right)^{(-1)} & =\Lambda_{2}^{(-1)} X_{k}^{(2)}\left(\left(X_{k}\right)^{(1)}\right)^{(-1)} \Lambda_{1}\left(X_{k+1}^{(1)}-\Lambda_{1}^{(-1)} \Delta_{k}^{(1)}\right)\left(X_{k+1}\right)^{(-1)} \\
& +\Lambda_{2}^{(-1)} \Delta_{k}^{(2)}\left(X_{k+1}\right)^{(-1)} \\
& =\Lambda_{2}^{(-1)} X_{k}^{(2)}\left(\left(X_{k}\right)^{(1)}\right)^{(-1)} \Lambda_{1} X_{k+1}^{(1)}\left(X_{k+1}\right)^{(-1)} \\
& -\Lambda_{2}^{(-1)} X_{k}^{(2)}\left(\left(X_{k}\right)^{(1)}\right)^{(-1)} \Lambda_{1} \Lambda_{1}^{(-1)} \Delta_{k}^{(1)}\left(X_{k+1}\right)^{(-1)}+\Lambda_{2}^{(-1)} \Delta_{k}^{(2)}\left(X_{k+1}\right)^{(-1)} \\
& \left.=\Lambda_{2}^{(-1)} X_{k}^{(2)}\left(\left(X_{k}\right)^{(1)}\right)^{(-1)} \Lambda_{1}-\Lambda_{2}^{(-1)} X_{k}^{(2)}\left(\left(X_{k}\right)^{(1)}\right)^{(-1)} \Delta_{k}^{(1)}\right)\left(X_{k+1}\right)^{(-1)} \\
& +\Lambda_{2}^{(-1)} \Delta_{k}^{(2)}\left(X_{k+1}\right)^{(-1)} \\
& =\Lambda_{2}^{(-1)} X_{k}^{(2)}\left(\left(X_{k}\right)^{(1)}\right)^{(-1)} \Lambda_{1} \\
& -\left(X_{k+1}\right)^{(-1)}\left(\Lambda_{2}^{(-1)} X_{k}^{(2)}\left(\left(X_{k}\right)^{(1)}\right)^{(-1)} \Delta_{k}^{(1)}-\Lambda_{2}^{(-1)} \Delta_{k}^{(2)}\right)
\end{aligned}
$$

and now let in the (5)

$$
\begin{aligned}
& =\Lambda_{2}^{(-1)} X_{k}^{(2)}\left(\left(X_{k}\right)^{(1)}\right)^{(-1)} \Lambda_{1}-\left(\Lambda_{2}^{(-1)} X_{k}^{(2)}\left(\left(X_{k}\right)^{(1)}\right)^{(-1)} \Delta_{k}^{(1)}\right. \\
& \left.-\Lambda_{2}^{(-1)} \Delta_{k}^{(2)}\right)\left(\Lambda_{1}^{(-1)} X_{k}^{(1)} R_{k+1}^{(-1)}+\Lambda_{1}^{(-1)} \Delta_{k}^{(1)}\right)^{-1}
\end{aligned}
$$

since $\Delta_{k}^{(i)}=U_{i}^{H} E_{k} R_{k+1}^{(-1)}$ and

$$
\begin{gathered}
\left(X_{k+1}^{(1)}\right)^{-1}=\left(\Lambda_{1}^{(-1)} X_{k}^{(1)} R_{k+1}^{(-1)}+\Lambda_{1}^{(-1)} \Delta_{k}^{(1)}\right)^{-1}=\left(\Lambda_{1}^{(-1)}\left(X_{k}^{(1)} R_{k+1}^{(-1)}+\Delta_{k}^{(1)}\right)\right)^{-1} \\
=\left(X_{k}^{(1)} R_{k+1}^{(-1)}\left(I+\Delta_{k}^{(1)} R_{k+1}\left(X_{k}^{(1)}\right)^{-1}\right)^{-1} \Lambda_{1}\right. \\
=R_{k+1}\left(X_{k}^{(1)}\right)^{-1}\left(I+\Delta_{k}^{(1)} R_{k+1}\left(X_{k}^{(1)}\right)^{-1}\right)^{-1} \Lambda_{1}
\end{gathered}
$$

Then we further simplify the expression $X_{k+1}^{(2)}\left(X_{k+1}^{(1)}\right)^{-1}$ to

$$
\begin{aligned}
X_{k+1}^{(2)}\left(X_{k+1}^{(1)}\right)^{-1}= & \Lambda_{2}^{-1} X_{k}^{(2)}\left(X_{k}^{(1)}\right)^{-1} \Lambda_{1} \\
& -\Lambda_{2}^{-1} X_{k}^{(2)}\left(X_{k}^{(1)}\right)^{-1} \Delta_{k}^{(1)}\left(X_{k+1}^{(1)}\right)^{-1}+\Lambda_{2}^{-1} \Delta_{k}^{(2)}\left(X_{k+1}^{(1)}\right)^{-1} \\
= & \Lambda_{2}^{-1} X_{k}^{(2)}\left(X_{k}^{(1)}\right)^{-1} \Lambda_{1}-\Lambda_{2}^{-1} X_{k}^{(2)}\left(X_{k}^{(1)}\right)^{-1} U_{1}^{H} E_{k} R_{k+1}^{-1}\left(X_{k+1}^{(1)}\right)^{-1} \\
& +\Lambda_{2}^{-1} U_{2}^{H} E_{k} R_{k+1}^{-1}\left(X_{k+1}^{(1)}\right)^{-1} \\
= & \Lambda_{2}^{-1} X_{k}^{(2)}\left(X_{k}^{(1)}\right)^{-1} \Lambda_{1}-\left(\Lambda_{2}^{-1} X_{k}^{(2)}\left(X_{k}^{(1)}\right)^{-1} U_{1}^{H}-\Lambda_{2}^{-1} U_{2}^{H}\right) E_{k} R_{k+1}^{-1}\left(X_{k+1}^{(1)}\right)^{-1} \\
\left(X_{k+1}^{(1)}\right)^{-1}= & \left(\Lambda_{1}^{-1} X_{k}^{(1)} R_{k+1}^{-1}+\Lambda_{1}^{-1} \Delta_{k}^{(1)}\right)^{-1}=\left(\Lambda_{1}^{-1} X_{k}^{(1)} R_{k+1}^{-1}+\Lambda_{1}^{-1} U_{1}^{H} E_{k} R_{k+1}^{-1}\right)^{-1} \\
= & \left(\Lambda_{1}^{-1}\left(X_{k}^{(1)}\right) R_{k+1}^{-1}\right)^{-1}\left(I+U_{1}^{H} E_{k}\left(X_{k}^{(1)}\right)^{-1}\right)^{-1} \\
= & R_{k+1}\left(X_{k}^{(1)}\right)^{-1}\left(I+U_{1}^{H} E_{k}\left(X_{k}^{(1)}\right)^{-1}\right)^{-1} \Lambda_{1} \\
X_{k+1}^{(2)}\left(X_{k+1}^{(1)}\right)^{(-1)}= & \Lambda_{2}^{-1} X_{k}^{(2)}\left(X_{k}^{(1)}\right)^{-1} \Lambda_{1} \\
& \quad\left(\Lambda_{2}^{-1} X_{k}^{(2)}\left(X_{k}^{(1)}\right)^{-1} U_{1}^{H}-\Lambda_{2}^{-1} U_{2}^{H}\right) E_{k}\left(X_{k}^{(1)}\right)^{-1}\left(I+U_{1}^{H} E_{k}\left(X_{k}^{(1)}\right)^{-1}\right)^{-1} \Lambda_{1}
\end{aligned}
$$

Taking 2-norm of the above equation at both sides and using the condition

$$
\begin{equation*}
\|V\|_{2}\|U\|_{2}\left(1+t_{k}+t_{k^{\prime}}\right) \epsilon_{k}<1, \tag{6}
\end{equation*}
$$

we obtain the following upper bound of $t_{k+1}$ :
$\left\|X_{k+1}^{(2)}\left(X_{k+1}^{(1)}\right)^{(-1)}\right\| \leqslant\left\|\Lambda_{2}^{-1}\right\|_{2}\left\|X_{k}^{(2)}\left(X_{k}^{(1)}\right)^{-1}\right\|_{2}\left\|\Lambda_{1}\right\|_{2}$

$$
-\left(\left\|\Lambda_{2}^{-1}\right\|_{2}\left\|X_{k}^{(2)}\left(X_{k}^{(1)}\right)^{-1}\right\|_{2}\left\|U_{1}^{H}\right\|_{2}\right.
$$

$\left.-\left\|\Lambda_{2}^{-1}\right\|_{2}\left\|U_{2}^{H}\right\|_{2}\right)\left\|E_{k}\right\|_{2}\left\|\left(X_{k}^{(1)}\right)^{-1}\right\|_{2}\left\|\left(I+U_{1}^{H} E_{k}\left(X_{k}^{(1)}\right)^{-1}\right)^{-1}\right\|_{2}\left\|\Lambda_{1}\right\|_{2}$

$$
\begin{aligned}
& =\left\|\Lambda_{2}^{-1}\right\|_{2}\left\|\Lambda_{1}\right\|_{2} t_{k} \\
& +\left(\left\|\Lambda_{2}^{-1}\right\|_{2} t_{k}\left\|U_{1}^{H}\right\|_{2}\right. \\
& \left.+\left\|\Lambda_{2}^{-1}\right\|_{2}\left\|U_{2}^{H}\right\|_{2}\right)\left\|E_{k}\right\|_{2}\left\|\left(X_{k}^{(1)}\right)^{-1}\right\|_{2}\left\|\left(I+U_{1}^{H} E_{k}\left(X_{k}^{(1)}\right)^{-1}\right)^{-1}\right\|_{2}\left\|\Lambda_{1}\right\|_{2}
\end{aligned}
$$

$$
\begin{aligned}
X_{k+1}^{(3)}\left(X_{k+1}^{(1)}\right)^{(-1)} & =\left\|\Lambda_{2}^{-1}\right\|_{2}\left\|\Lambda_{1}\right\|_{2} t_{k^{\prime}} \\
& +\left(\left\|\Lambda_{2}^{-1}\right\|_{2} t_{k^{\prime}}\left\|U_{1}^{H}\right\|_{2}\right. \\
& \left.+\left\|\Lambda_{2}^{-1}\right\|_{2}\left\|U_{2}^{H}\right\|_{2}\right)\left\|E_{k}\right\|_{2}\left\|\left(X_{k}^{(1)}\right)^{-1}\right\|_{2}\left\|\left(I+U_{1}^{H} E_{k}\left(X_{k}^{(1)}\right)^{-1}\right)^{-1}\right\|_{2}\left\|\Lambda_{1}\right\|_{2} \\
t_{k+1} \leqslant \rho\left(t_{k}+\right. & \left.t_{k^{\prime}}\right)+\left(\left(\left\|\Lambda_{2}^{-1}\right\|_{2} t_{k}+\left\|\Lambda_{2}^{-1}\right\|_{2}+\left\|\Lambda_{2}^{-1}\right\|_{2} t_{k^{\prime}}+\left\|\Lambda_{2}^{-1}\right\|_{2}\right)\right. \\
& \times\|U\|_{2}\left\|E_{k}\right\|_{2}\left\|\left(X_{k}^{(1)}\right)^{-1}\right\|_{2}\left\|\Lambda_{1}\right\|_{2}\left\|\left(I+U_{1}^{H} E_{k}\left(X_{k}^{(1)}\right)^{-1}\right)^{-1}\right\|_{2} \\
\leqslant \rho\left(t_{k}+\right. & \left.t_{k^{\prime}}\right)+\left(1+t_{k}+t_{k^{\prime}}\right)\left\|\Lambda_{1}^{-1}\right\|_{2} \frac{\|U\|_{2}\left\|E_{k}\right\|_{2}\left\|\left(X_{k}^{(1)}\right)^{-1}\right\|_{2}\left\|\Lambda_{1}\right\|_{2}}{1-\left\|E_{k}\right\|_{2}\left\|\left(X_{k}^{(1)}\right)^{-1}\right\|_{2}\|U\|_{2}}
\end{aligned}
$$

Using lemma (3.3), we know that $X_{k}^{(1)} \leqslant\|v\|_{2}\left(1+t_{k}+t_{k^{\prime}}\right)$. From this and (3), the final bound for $t_{k+1}$ is derived.

Lemma 3.5
Assume that $X_{0}$ is such that $X_{0}^{(1)}$ is invertible. If

$$
\epsilon_{k} \leqslant \epsilon:=\frac{(1-\rho) 2 t_{0}}{\|V\|_{2}\|U\|_{2}\left(1+2 t_{0}\right)\left(\rho+2 t_{0}\right)}
$$

for all $k$, then $t_{k} \leqslant t_{0}$.
Proof We prove $t_{k} \leqslant t_{0}, t_{k^{\prime}} \leqslant t_{0}$ by induction. Supposing $X_{k}^{(1)}$ is nonsingular and $t_{k} \leqslant t_{0}, t_{k^{\prime}} \leqslant t_{0}$ is true for some $k, k^{\prime}$, we show that $X_{k+1}^{(1)}$ is nonsingular and $t_{k+1} \leqslant 2 t_{0}$. First note that from $\epsilon_{k} \leqslant \epsilon$, we have

$$
\|V\|_{2}\|U\|_{2}\left(1+t_{k}+t_{k^{\prime}}\right) \epsilon_{k} \leqslant\|V\|_{2}\|U\|_{2}\left(1+t_{0}+t_{k^{\prime}}\right) \epsilon=\frac{(1+\rho) 2 t_{0}}{\rho+2 t_{0}}<1
$$

We discuss in two cases:

- Case I
$X_{k+1}^{(1)}$ is nonsingular. Then by lemma (3.4), we have

$$
\begin{aligned}
t_{k+1} & \leqslant \rho\left(t_{k}+t_{k^{\prime}}\right)+\frac{\rho\|V\|_{2}\|U\|_{2}\left(1+t_{k}+t_{k^{\prime}}\right)^{2} \epsilon_{k}}{1-\|V\|_{2}\|U\|_{2}\left(1+t_{k}+t_{k^{\prime}}\right) \epsilon_{k}} \\
& \leqslant \rho 2 t_{0}+\frac{\rho\|V\|_{2}\|U\|_{2}\left(1+2 t_{0}\right)^{2} \epsilon}{1-\|V\|_{2}\|U\|_{2}\left(1+2 t_{0}\right) \epsilon} \\
& \leqslant \rho 2 t_{0}+\frac{\rho\left(1+2 t_{0}\right) \frac{(1-\rho) 2 t_{0}}{\rho+2 t_{0}}}{1-\frac{(1-\rho) 2 t_{0}}{\rho+2 t_{0}}} \leqslant \rho 2 t_{0}+\frac{\frac{\rho\left(1+2 t_{0}\right)(1-\rho) 2 t_{0}}{\rho+2 t_{0}}}{\frac{\rho+2 t_{0}+2 t_{0}+\rho 2 t_{0}}{\rho+2 t_{0}}} \\
& \leqslant \rho 2 t_{0}+\frac{\rho\left(1+2 t_{0}\right)(1-\rho) 2 t_{0}}{\rho\left(1+2 t_{0}\right)}=2 t_{0},
\end{aligned}
$$

- Case II
$X_{k+1}^{(1)}$ is singular. Then let

$$
\tilde{y}_{k+1}=y_{k+1}+\delta V_{1} R_{k+1}+\mu V_{2}\binom{1}{0} R_{k+1}+\theta V_{3}\binom{0}{1} R_{k+1}
$$

where $Y_{k+1}=X_{k+1} R_{k+1}$ and $\delta, \mu>0$ are two parameters. Then we have

$$
\begin{aligned}
& A y_{k+1}^{\tilde{k}}=A y_{k+1}+\delta A V_{1} R_{k+1}+\mu A V_{2}\binom{1}{0} R_{k+1}+\theta V_{3}\binom{0}{1} R_{k+1} \\
& A y_{k+1}^{\tilde{k}}=B X_{k}+E_{k}+\delta A V_{1} R_{k+1}+\mu A V_{2}\binom{1}{0} R_{k+1}+\theta V_{3}\binom{0}{1} R_{k+1} \\
& E_{k}+\delta A V_{1} R_{k+1}+\mu A V_{2}\binom{1}{0} R_{k+1}+\theta V_{3}\binom{0}{1} R_{k+1}=\tilde{E}_{k} \\
& A y_{\tilde{k+1}}=B X_{k}+\tilde{E}_{k}
\end{aligned}
$$

Since $\left\|E_{k}\right\|_{2} \leqslant \epsilon_{k}$, we have $\left\|\tilde{E}_{k}\right\|_{2} \leqslant \epsilon_{k}$ for sufficiently small $\delta$ and $\mu$. Let $\tilde{Y}_{k+1}=$ $\tilde{X}_{k+1} \tilde{R}_{k+1}$ be the QR-factorization and let $\tilde{X}_{k+1}=V_{1} \tilde{X}_{(k+1)}^{(1)}+V_{2} \tilde{X}_{k+1}^{(2)}+V_{3} \tilde{X}_{k+1}^{(3)}$. Then $\tilde{X}_{k+1}$ satisfies the same condition that $X_{k+1}$ does and the bound on $t_{k+1}$ applies to $\tilde{t}_{k+1}:=\left\|\tilde{X}_{k+1}^{(2)}\left(\tilde{X}_{k+1}^{(1)}\right)^{-1}\right\|_{2}$ as well. It follows from

$$
\begin{gather*}
\tilde{Y}_{k+1}=V_{1} X_{k+1}^{(1)} \tilde{R}_{k+1}+V_{2} \tilde{X}_{k+1}^{(2)} \tilde{R}_{k+1}+V_{3} \tilde{X}_{k+1}^{(3)} \tilde{R}_{k+1}  \tag{7}\\
\tilde{Y}_{k+1}=Y_{k+1}+\delta V_{1} R_{k+1}+\mu V_{2}\binom{1}{0} R_{k+1}+\theta V_{3} R_{k+1}\binom{0}{1} \\
=\left(Y_{k+1} R_{k+1}^{-1}+\delta V_{1}+\mu V_{2}\binom{1}{0}+\theta V_{3}\binom{0}{1}\right) R_{k+1} \\
=\left(V_{1} X_{k+1}^{(1)}+V_{2} X_{k+1}^{(2)}+\delta V_{1}+\mu V_{2}\binom{1}{0}+\theta V_{3}\binom{0}{1}\right) R_{k+1} \\
\left.=\left[V_{1}\left(X_{k+1}^{(1)}+\delta I\right)+V_{2}\left(X_{k+1}^{(2)}+\mu\binom{1}{0}\right)+V_{3}\left(X_{k+1}^{(3)}+\theta\binom{0}{1}\right)\right)\right] R_{k+1} \tag{8}
\end{gather*}
$$

now let this relation (7) is equal to (8) that

$$
\begin{aligned}
& X_{k+1}^{(1)}=\left(X_{k+1}^{(1)}+\delta I\right) R_{k+1} \tilde{R}_{k+1}^{(-1)} \\
& X_{k+1}^{(2)}=\left(X_{k+1}^{(2)}+\mu\binom{1}{0}\right) R_{k+1} \tilde{R}_{k+1}^{(-1)} \quad X_{k+1}^{(3)}=\left(X_{k+1}^{(3)}+\theta\binom{0}{1}\right) R_{k+1} \tilde{R}_{k+1}^{(-1)}
\end{aligned}
$$

So $X_{k+1}^{(1)}$ is nonsingular for sufficiently small $\delta>0$. Then, by case I , we have

$$
\tilde{t}_{k+1}=\left\|\tilde{X}_{k+1}^{(2)}\left(\tilde{X}_{k+1}^{(1)}\right)^{(-1)}\right\|_{2} \leqslant t_{0}
$$

$$
\tilde{t}_{k+1}=\left\|\tilde{X}_{k+1}^{(3)}\left(\tilde{X}_{k+1}^{(1)}\right)^{(-1)}\right\|_{2} \leqslant t_{0}
$$

for all sufficiently small $\delta>0$ and $\mu \leqslant 0$ and $\theta \leqslant 0$. However

$$
\begin{aligned}
\tilde{X}_{k+1}^{(2)}\left(\tilde{X}_{k+1}^{(1)}\right)^{(-1)} & =\left(X_{k+1}^{(2)}+\mu\binom{1}{0}\right) R_{k+1} \tilde{R}_{k+1}^{(-1)}\left(X_{k+1}^{(1)}+\delta I\right)^{(-1)} \tilde{R}_{k+1} R_{k+1}^{(-1)} \\
& =\left(X_{k+1}^{(2)}+\mu\binom{1}{0}\right)\left(X_{k+1}^{(1)}+\delta I\right)^{(-1)} \\
& =X_{k+1}^{(2)}\left(X_{k+1}^{(1)}+\delta I\right)^{(-1)}+\mu\binom{1}{0}\left(\left(X_{k+1}^{(1)}+\delta I\right)^{(-1)}\right)
\end{aligned}
$$

is unbounded as $\delta \rightarrow 0$, because if $X_{k+1}^{(2)}\left(X_{k+1}^{(1)}+\delta I\right)^{-1}$ is unbounded, then $\tilde{t}_{k+1}$ is unbounded by setting $\mu=0$; and if $X_{k+1}^{(2)}\left(X_{k+1}^{(1)}+\delta I\right)^{-1}$ is bounded, then $\tilde{t}_{k+1}$ is unbounded by setting $\mu>0$. Therefore $X_{k+1}^{(1)}$ is nonsingular and hence $t_{k+1} \leqslant 2 t_{0}$

$$
\begin{aligned}
\tilde{X}_{k+1}^{(3)}\left(\tilde{X}_{k+1}^{(1)}\right)^{(-1)} & =\left(X_{k+1}^{(3)}+\theta\binom{1}{0}\right) R_{k+1} \tilde{R}_{k+1}^{(-1)}\left(X_{k+1}^{(1)}+\delta I\right)^{(-1)} \tilde{R}_{k+1} R_{k+1}^{(-1)} \\
& =\left(X_{k+1}^{(3)}+\theta\binom{0}{1}\right)\left(X_{k+1}^{(1)}+\delta I\right)^{(-1)} \\
& =X_{k+1}^{(3)}\left(X_{k+1}^{(1)}+\delta I\right)^{(-1)}+\theta\binom{0}{1}\left(\left(X_{k+1}^{(1)}+\delta I\right)^{(-1)}\right)
\end{aligned}
$$

is unbounded as $\delta \rightarrow 0$, because if $X_{k+1}^{(3)}\left(X_{k+1}^{(1)}+\delta I\right)^{-1}$ is unbounded, then $\tilde{t}_{k+1}$ is unbounded by setting $\theta=0$; and if $X_{k+1}^{(3)}\left(X_{k+1}^{(1)}+\delta I\right)^{-1}$ is bounded, then $\tilde{t}_{k+1}$ is unbounded by setting $\theta>0$. Therefore $X_{k+1}^{(1)}$ is nonsingular and hence $t_{k+1} \leqslant 2 t_{0}$ proof is completed.

We now prove our main result on convergence of $t_{k}$ and $t_{k^{\prime}}$. We are interested in the case that $\epsilon_{k}$ is a linearly decreasing sequence.

Theorem 3.6 Assume that $X_{0}$ is such that $X_{0}^{-1}$ is invertible. Let $\epsilon_{k}=a \gamma^{k}$ with $\gamma<1$ and

$$
a \leqslant \frac{(1-\rho) 2 t_{0}}{\|V\|_{2}\|U\|_{2}\left(1+2 t_{0}\right)\left(\rho+2 t_{0}\right)}
$$

Then we have

$$
t_{k} \leqslant\left\{\begin{array}{l}
2 \rho^{k} t_{0}+a c \frac{\gamma^{k}-\rho^{k}}{\frac{-\rho}{\gamma}} \gamma \neq \rho \\
2 \rho^{k} t_{0}+a c k \rho^{k-1} \\
\gamma=\rho
\end{array}\right.
$$

where

$$
c=\|V\|_{2}\|U\|_{2}\left(1+2 t_{0}\right)\left(\rho+2 t_{0}\right)
$$

Proof Since

$$
\epsilon_{k} \leqslant \frac{(1-\rho) 2 t_{0}}{\|V\|_{2}\|U\|_{2}\left(1+2 t_{0}\right)\left(\rho+2 t_{0}\right)}
$$

we have $t_{k} \leqslant t_{0}$ by Lemma (3.5). Then,

$$
\epsilon_{k} \leqslant \frac{(1-\rho) 2 t_{0}}{\|V\|_{2}\|U\|_{2}\left(1+2 t_{0}\right)\left(\rho+2 t_{0}\right)}
$$

so we have $t_{k} \leqslant t_{0}$

$$
\|V\|_{2}\|U\|_{2}\left(1+t_{k}+t_{k^{\prime}}\right) \epsilon_{k} \leqslant\|V\|_{2}\|U\|_{2}\left(1+2 t_{0}\right) \epsilon \leqslant \frac{(1-\rho) 2 t_{0}}{\left(\rho+2 t_{0}\right)}<1
$$

It follows from Lemma (3.4) that $t_{k+1} \leqslant \rho\left(t_{k}+t_{k^{\prime}}\right)+c_{k} \epsilon_{k}$

$$
\begin{aligned}
t_{k+1} & \leqslant \rho\left(t_{k}+t_{k^{\prime}}\right)+\frac{\rho\|V\|_{2}\|U\|_{2}\left(1+t_{k}+t_{k^{\prime}}\right)^{2} \epsilon_{k}}{1-\|V\|_{2}\|U\|_{2}\left(1+t_{k}+t_{k^{\prime}}\right) \epsilon_{k}} \\
t_{k+1} & \leqslant \rho\left(t_{k}+t_{k^{\prime}}\right)+a c_{k} \gamma^{k} \\
t_{k} & \leqslant \rho\left(t_{k}+t_{k^{\prime}}\right)+a c_{k-1} \gamma^{k-1} \\
\quad & \\
t_{k} & \leqslant \rho^{k} 2 t_{0}+a c \frac{\gamma^{k}-\rho^{k}}{\gamma-\rho}
\end{aligned}
$$

now if $\rho=\gamma$ we have $t_{k} \leqslant \rho^{k} 2 t_{0}+a c k \rho^{k-1}$.

$$
\begin{aligned}
t_{k+1} & \leqslant \rho\left(t_{k}+t_{k^{\prime}}\right)+c_{k} \epsilon_{k} \\
c_{k} & =\frac{\rho\|V\|_{2}\|U\|_{2}\left(t_{k}+t_{k_{k}}\right)^{2}}{1-\|V\|_{2}\|U\|_{2}\left(t_{k}+t_{k^{\prime}}\right) \epsilon_{k}} \leqslant \frac{\rho\|V\|_{2}\|U\|_{2}\left(1+2 t_{0}\right)^{2}}{1-\frac{(1-\rho) 2 t_{0}}{\left(\rho+2 t_{0}\right)}} \\
& \leqslant \frac{\rho\|V\|_{2}\|U\|_{2}\left(1+2 t_{0}\right)^{2}}{\frac{\rho+t_{0}-t_{0}+\rho t_{0}}{\left(\rho+2 t_{0}\right)}} \leqslant \frac{\rho\|V\|_{2}\|U\|_{2}\left(1+2 t_{0}\right)^{2}\left(\rho+2 t_{0}\right)}{\rho\left(1+2 t_{0}\right)^{2}} \\
& \leqslant\|V\|_{2}\|U\|_{2}\left(1+2 t_{0}\right)\left(\rho+2 t_{0}\right)=c
\end{aligned}
$$

Therefore, $t_{k+1} \leqslant \rho\left(t_{k}+t_{k^{\prime}}\right)+a c \gamma^{k}$. Solving this inequality, we obtain the bound for $t_{k}$

The conclusion of the above theorem is that the subspace spanned by $X_{k}, R\left(X_{k}\right)$, converges to the spectral subspace $R\left(V_{1}\right)$ linearly at the rate of $\max \{\rho, \gamma\}$. The condition on a is to ensure convergence and is clearly not a necessary condition. An interesting fact is that there is no gain in convergence rate if we choose $\gamma<\rho$, some shall focus on the case $\gamma>\rho$. The following corollary gives a more precise bound for the constant C and hence for $t_{k}$ and $t_{k^{\prime}}$ at the convergence stage.

Corollary 3.7
Let $1>\gamma>\rho$ and $\epsilon_{k}=a^{k}$. Suppose that $a$ is chosen such that $t_{k} \longrightarrow 0$ and
$t_{k^{\prime}} \longrightarrow 0$. Then

$$
\lim \sup \frac{t_{k}}{a \gamma^{k}} \leqslant \rho(\gamma-\rho)^{-1}\|V\|_{2}\|U\|_{2}
$$

and

$$
\lim \sup \frac{t_{k^{\prime} 0}}{a \gamma^{k}} \leqslant \rho(\gamma-\rho)^{-1}\|V\|_{2}\|U\|_{2}
$$

Proof Apply the main theorem to $t_{k}$ starting from $k=k_{0}$, we have

$$
t_{k} \leqslant 2 \rho^{k-k_{0}} t_{k_{0}}+\frac{\gamma^{k-k_{0}}-\rho^{k-k_{0}}}{\gamma-\rho} a \gamma^{k_{0}} c_{k_{0}}
$$

where

$$
\begin{gathered}
c_{k_{0}}=\frac{\rho\|V\|_{2}\|U\|_{2}\left(1+t_{k_{0}}+t_{k^{\prime} 0}\right)^{2}}{1-\rho\|V\|_{2}\|U\|_{2}\left(1+t_{k_{0}}+t_{k^{\prime}}\right) \epsilon_{k_{0}}} \\
c_{k_{0}}=\frac{\rho\|V\|_{2}\|U\|_{2}\left(1+t_{k_{0}}+t_{k^{\prime}}\right)^{2}}{1-\rho\|V\|_{2}\|U\|_{2}\left(1+t_{k_{0}}+t_{k^{\prime} 0}\right) \epsilon_{k_{0}}} \sim \rho\|V\|_{2}\|U\|_{2} \\
\|V\|_{2}\|U\|_{2}\left(1+2 t_{k_{0}}\right)\left(\rho+2 t_{k_{0}}\right) \leqslant \frac{2\left(1-\rho_{k_{0}}\right) t_{k_{0}}}{\epsilon_{k_{0}}}
\end{gathered}
$$

Dividing $a \gamma^{k}$ and taking $k \longrightarrow \infty$ first and then $k_{0} \longrightarrow \infty$ in the inequality, we obtain the bound.

$$
\begin{aligned}
& \frac{t_{k}}{a \gamma^{k}} \leqslant \frac{\rho^{k-k_{0}} t_{k_{0}}}{a \gamma^{k}}+\frac{\gamma^{k-k_{0}}-\rho^{k-k_{0}}}{(\gamma-\rho) a \gamma^{k-k_{0}}} a \gamma^{k_{0}} c_{k_{0}} \\
& \frac{t_{k}}{a \gamma^{k}} \leqslant \frac{\rho^{k-k_{0}} t_{k_{0}}}{a \gamma^{k}}+\frac{\gamma^{k-k_{0}}-\rho^{k-k_{0}}}{(\gamma-\rho)} \gamma^{-\left(k-k_{0}\right)} c_{k_{0}}
\end{aligned}
$$

and so $\rho<\gamma<1$

$$
\lim _{k \rightarrow \infty}\left(\frac{\rho}{\gamma}\right)^{k}\left(\frac{\rho^{-k_{0}} t_{k_{0}}}{a}\right)=0
$$

and

$$
\lim _{k \rightarrow \infty}\left(\frac{\rho}{\gamma}\right)\left(\frac{\rho^{-k_{0}} t_{k_{0}}}{a}\right)+\lim _{k \rightarrow \infty} \frac{\gamma^{k-k_{0}}\left(1-\left(\frac{\rho}{\gamma}\right)^{k-k_{0}}\right)}{(\gamma-\rho)} \gamma^{-\left(k-k_{0}\right)} c_{k_{0}}
$$

and

$$
\begin{gathered}
=0+\lim _{k \rightarrow \infty} \frac{1-0}{\gamma-\rho}\left(\rho\|V\|_{2}\|U\|_{2}\right)=\rho(\gamma-\rho)^{-1}\|V\|_{2}\|U\|_{2} \\
\lim \sup \frac{t_{k}}{a \gamma^{k}} \leqslant \rho(\gamma-\rho)^{-1}\|V\|_{2}\|U\|_{2}
\end{gathered}
$$

Apply the main theorem to $t_{k^{\prime}{ }_{0}}$ starting from $k^{\prime}=k_{0}$, we have

$$
\limsup \frac{t_{k^{\prime}}}{a \gamma^{k}} \leqslant \rho(\gamma-\rho)^{-1}\|V\|_{2}\|U\|_{2}
$$

## 4. Conclusions

We have presented an inexact inverse subspace iteration for computing a few smallest eigenpairs of the generalized eigenvalue problem $A x=B x$. By properly scaling the block vectors, we ensure convergence of columns in the iterative blocks, which allows using approximation from one step as an initial approximation for the next step.we analyzed convergence of the subspace to the spectral space sought.

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